

CONGRUENCES AND IDEALS ON BOOLEAN MODULES: A HETEROGENEOUS POINT OF VIEW

SANDRA MARQUES PINTO AND M. TERESA OLIVEIRA-MARTINS

ABSTRACT: Definitions for heterogeneous congruences and heterogeneous ideals on a Boolean module \mathcal{M} are given and the respective lattices $Cong\mathcal{M}$ and $Ide\mathcal{M}$ are presented. A characterization of the simple Boolean modules is achieved differing from that given by Brink in a homogeneous approach. We construct the smallest and the greatest modular congruence having the same Boolean part. The same is established for modular ideals. The notions of kernel of a modular congruence and the congruence induced by a modular ideal are introduced to describe an isomorphism between $Cong\mathcal{M}$ and $Ide\mathcal{M}$. This isomorphism leads us to conclude that the class of the Boolean module is ideal determined.

KEYWORDS: Relation algebras; Boolean modules; modular heterogeneous congruence; modular heterogeneous ideal; simple Boolean module.

AMS SUBJECT CLASSIFICATION (2010): 03B70, 03G05, 03G15, 06B10, 06E25, 08A68.

1. Introduction

The application of abstract algebra in logic and computer science rely and depends on the simultaneous study of algebras of sets and algebras of binary relations. To talk about algebras of sets is synonymous to study Boolean algebras and the most famous algebra of relations is that presented by Tarski in [9]. There Tarski introduces the relations algebras as algebras of binary relations adding to the Boolean structure the operations of composition, converse and identity. Boolean modules were first established by Brink in [1]. Given a relation algebra \mathcal{R} , Brink defined and studied Boolean \mathcal{R} -modules as a Boolean algebra \mathcal{B} with actions from the relation algebra \mathcal{R} . Such actions were induced by a map called Peircean operator, $:$, from $\mathcal{R} \times \mathcal{B}$ to \mathcal{B} , where each element $a \in \mathcal{R}$ defines in \mathcal{B} a map

$$\begin{aligned} \mathcal{B} &\rightarrow \mathcal{B} \\ p &\mapsto a : p \end{aligned}$$

Received July 29, 2010.

Financial support from the *CMUC, Department of Mathematics, University of Coimbra, 3001-454 Coimbra, Portugal*, is gratefully acknowledged.

satisfying a required set of axioms. A unified concept associated to this homogeneous approach is given naturally by a two sorted algebra $\mathcal{M} = (\mathcal{B}, \mathcal{R}, :)$ containing a Boolean algebra \mathcal{B} and a relation algebra \mathcal{R} where the Peircean operator, $:$, is interpreted now as a heterogeneous operation in \mathcal{M} ranging from $\mathcal{R} \times \mathcal{B}$ to \mathcal{B} . The importance of the heterogeneous algebras approach on Boolean modules is fully presented on the introduction of [6] by R. Hirsch and [2] by Brink, Britz and Schmidt. Nevertheless, their characterization of simple Boolean modules follows a homogeneous point of view, since their definition of a Boolean module ideal is a Boolean algebra ideal closed under multiplication by elements of the relation algebra. The same can be stated concerning congruences. A throughout heterogeneous approach is followed in our work for both concepts under study (Definitions 3.1 and 4.1). Thus a modular congruence θ is considered as adequate pair of congruences $\theta = (\theta_1, \theta_2)$ with θ_1 a Boolean congruence and θ_2 a relation congruence and modular ideals I as suitable pairs of ideals $I = (I_1, I_2)$ where I_1 is a Boolean ideal and I_2 is a ideal on the relation algebra.

2. Boolean modules

Boolean modules were introduced by Brink [1] as homogeneous algebras, Boolean algebras with a multiplication (Peircean product) from a relation algebra. A Boolean module is, from a heterogeneous point of view, a two sorted algebra containing a Boolean algebra, a relation algebra and an operator (a heterogeneous operation, the Peircean operator) taking a pair of a relation algebra element and a Boolean algebra element and originating a Boolean algebra element. We present here the standard definition of relation algebras given by Brink (originated from Chin and Tarski [3] and modified in Tarski [10]).

Definition 2.1. A *relation algebra* is an algebra $\mathcal{R} = (R, \vee, \wedge, ', \circ, 1, ;, \checkmark, e)$ satisfying the following axioms for each $a, b, c \in R$

R1 $(R, \vee, \wedge, ', \circ, 1)$ is a Boolean algebra

R2 $a; (b; c) = (a; b); c$

R3 $a; e = a = e; a$

R4 $a^{\checkmark\checkmark} = a$

R5 $(a \vee b); c = a; c \vee b; c$

R6 $(a \vee b)^{\checkmark} = a^{\checkmark} \vee b^{\checkmark}$

R7 $(a; b)^{\checkmark} = b^{\checkmark}; a^{\checkmark}$

R8 $a^{\checkmark}; (a; b)' \leq b'$.

Notation. For $a, b \in R$ we also write ab instead of $a; b$.

As usual, for elements p, q on a Boolean algebra B we define $p \oplus q = (p \wedge q') \vee (p' \wedge q)$. In particular, for elements a, b on a relation algebra R we also define $a \oplus b = (a \wedge b') \vee (a' \wedge b)$.

The standard class of models of relation algebras is the class of proper relation algebras.

Definition 2.2. A *proper relation algebra* over a non-empty set U is a set of binary relations on U that contains the identity relation and is closed with respect to union, intersection, complementation, relational composition and converse. If a proper relation algebra consists of all binary relations defined on U , then this algebra is called the *full relation algebra* and is denoted by $\mathcal{R}(U)$. More precisely, $\mathcal{R}(U)$ is the power set algebra over U^2 endowed with composition (“;”), converse (“ \checkmark ”) and identity (“ Id ”) operations defined, for $a, b \subseteq U^2$, by

$$a; b = \{(s, t) : \text{exists } u \in U \text{ such that } (s, u) \in a \text{ and } (u, t) \in b\}$$

$$a^\checkmark = \{(s, t) : (t, s) \in a\}$$

$$Id = \{(s, s) : s \in U\}.$$

The arithmetic of relation algebras can be described by the facts assembled on the following theorem.

Theorem 2.3. *On any relation algebra \mathcal{R} the following hold for any $a, b, c, d \in R$*

$$R9 \quad e^\checkmark = e, \quad o^\checkmark = o, \quad 1^\checkmark = 1$$

$$R10 \quad a \leq b \quad \text{if and only if} \quad a^\checkmark \leq b^\checkmark$$

$$R11 \quad (a \wedge b)^\checkmark = a^\checkmark \wedge b^\checkmark, \quad a'^\checkmark = a^{\checkmark\prime}$$

$$R12 \quad a; o = o = o; a, \quad 1; 1 = 1$$

$$R13 \quad a(b \vee c) = ab \vee ac$$

$$R14 \quad \text{If } a \leq b \text{ then } ca \leq cb \text{ and } ac \leq bc.$$

$$R15 \quad (ab) \wedge c = o \quad \text{if and only if} \quad (a^\checkmark c) \wedge b = o \quad \text{if and only if} \quad (cb)^\checkmark \wedge a = o$$

$$R16 \quad (ab) \wedge (cd) \leq a[(a^\checkmark c) \wedge (bd^\checkmark)]d$$

$$R17 \quad (a \oplus b)^\checkmark = a^\checkmark \oplus b^\checkmark.$$

Proof: R9-R16 are proved in [3]. To prove R17 we use R6 and R11. Thus $(a \oplus b)^\checkmark = [(a \wedge b') \vee (a' \wedge b)]^\checkmark = (a \wedge b')^\checkmark \vee (a' \wedge b)^\checkmark = (a^\checkmark \wedge b'^\checkmark) \vee (a'^\checkmark \wedge b^\checkmark) = (a^\checkmark \wedge b'^\checkmark) \vee (a'^\checkmark \wedge b^\checkmark) = a^\checkmark \oplus b^\checkmark.$ ■

As mentioned before, Brink introduced the notion of a Boolean \mathcal{R} -module \mathcal{B} as a homogeneous algebra. Here, the heterogeneous approach followed in

this study is emphasized from the very beginning, on the following definition, where the roles of \mathcal{B} and \mathcal{R} are taken evenly.

Definition 2.4. A *Boolean module* is a two-sorted algebra $\mathcal{M} = (\mathcal{B}, \mathcal{R}, :)$ where \mathcal{B} is a Boolean algebra, \mathcal{R} is a relation algebra and $:$ is a mapping $\mathcal{R} \times \mathcal{B} \longrightarrow \mathcal{B}$ (written $a : p$) such that for any $a, b \in R$ and $p, q \in B$, the following assertions are satisfied.

- M1 $a : (p \vee q) = a : p \vee a : q$
- M2 $(a \vee b) : p = a : p \vee b : p$
- M3 $a : (b : p) = (a; b) : p$
- M4 $e : p = p$
- M5 $o : p = 0$
- M6 $a^\checkmark : (a : p)' \leq p'$

Notation. For $a, b \in R$ and $p \in B$ we also use ap to represent $a : p$.

The standard models of Boolean modules are provided by the class of proper Boolean modules.

Definition 2.5. A *proper Boolean module* is a two-sorted algebra of a proper Boolean algebra (a field of sets) and a proper relation algebra together with *Peirce product* defined on sets and relations. For any relation a over some non-empty set U and any subset p of U , the *Peirce product* $:$ of a and p is defined by

$$a : p = \{s \in U : \text{there exists } t \in p \text{ such that } (s, t) \in a\}.$$

A *full Boolean module* $\mathcal{M}(U)$ over a non-empty set U is the Boolean module $(\mathcal{B}(U), \mathcal{R}(U), :)$, where $\mathcal{B}(U)$ is the power set algebra over U , $\mathcal{R}(U)$ is the full relation algebra over U and $:$ is the *Peirce product* defined set-theoretically.

Some facts valid on Boolean modules deserve mention.

Theorem 2.6. *On any Boolean module $\mathcal{M} = (\mathcal{B}, \mathcal{R}, :)$ the following hold for any $a, b \in R$ and $p, q \in B$*

- M7 *If $p \leq q$ then $ap \leq aq$.*
- M8 *If $a \leq b$ then $ap \leq bp$.*
- M9 $a(p \wedge q) \leq (ap \wedge aq)$
- M10 $(a \wedge b)p \leq (ap \wedge bp)$
- M11 $ap \wedge aq = 0$ *if and only if* $a^\checkmark q \wedge p = 0$
- M12 *If $\sum_{i \in I} p_i$ exists, then so does $\sum_{i \in I} ap_i$, and $a \sum_{i \in I} p_i = \sum_{i \in I} ap_i$.*
- M13 $a0 = 0$

- M14 $1 : 1 = 1$
M15 $(a1)' \leq a'1$
M16 $ap \wedge q \leq a(p \wedge a\check{q})$
M17 $p \leq 1p$

Proof: Proved in [1]. ■

3. The lattice $Cong\mathcal{M}$ of modular congruences

The concept of congruence with its recognized unifier formulation plays a central role both on lattice and universal algebra theories in general. Once more the presentation of next notion follows a heterogeneous view-point.

Definition 3.1. Let $\mathcal{M} = (\mathcal{B}, \mathcal{R}, :)$ be a Boolean module. A pair $\theta = (\theta_1, \theta_2)$ is a (*modular*) *congruence relation* on \mathcal{M} if θ_1 is a congruence relation on \mathcal{B} , θ_2 is a congruence relation on \mathcal{R} and $ap \theta_1 bq$ whenever $(p \theta_1 q$ and $a \theta_2 b)$.

Let us denote by $Cong\mathcal{M}$ the set of all modular congruences defined on a Boolean module \mathcal{M} . The set $Cong\mathcal{M}$ is partially ordered by $(\theta_1, \theta_2) \leq (\gamma_1, \gamma_2)$ if and only if $\theta_1 \subseteq \gamma_1$ and $\theta_2 \subseteq \gamma_2$. Our next aim is to define the lattice structure $(Cong\mathcal{M}, \wedge_{\mathcal{M}}, \vee_{\mathcal{M}})$. Since the intersection $\theta \cap \gamma = (\theta_1 \cap \gamma_1, \theta_2 \cap \gamma_2)$ of any two modular congruences θ and γ defined on \mathcal{M} is, itself, a modular congruence on \mathcal{M} , let $\theta \wedge_{\mathcal{M}} \gamma = \theta \cap \gamma$. Let us use $\langle \theta \rangle_{\mathcal{A}}$ to represent the congruence relation generated by the binary relation θ on any (homogeneous or heterogeneous) algebra \mathcal{A} , i.e., the intersection of all congruence relations θ' on \mathcal{A} containing θ

$$\langle \theta \rangle_{\mathcal{A}} = \cap \{ \theta' : \theta' \in Cong\mathcal{A} \text{ and } \theta \subseteq \theta' \}.$$

Now we need to define $\theta \vee_{\mathcal{M}} \gamma = (\tau_1, \tau_2)$. Using the classic definition of supremum of two congruences, the relation part of the congruence $\theta \vee_{\mathcal{M}} \gamma$ can be given by $\tau_2 = \theta_2 \vee_{\mathcal{R}} \gamma_2 = \langle \theta_2 \cup \gamma_2 \rangle_{\mathcal{R}}$. As far as the Boolean part is concerned some caution is required. Since the Boolean part must be closed to the operation $:$ evolving elements of R , we could be led to think of enlarging $\langle \theta_1 \cup \gamma_1 \rangle_{\mathcal{B}}$ with, for instance, elements of the type (ap, bq) with $(a, b) \in \theta_2$ and $(p, q) \in \gamma_1$. In fact, that is not necessary, as shown below.

Proposition 3.2. Let $\mathcal{M} = (\mathcal{B}, \mathcal{R}, :)$ be a Boolean module. For $\theta, \gamma \in Cong\mathcal{M}$, $(p, q) \in \langle \theta_1 \cup \gamma_1 \rangle_{\mathcal{B}}$ and $(a, b) \in \langle \theta_2 \cup \gamma_2 \rangle_{\mathcal{R}}$ we have $(ap, bq) \in \langle \theta_1 \cup \gamma_1 \rangle_{\mathcal{B}}$.

Proof: Analogous to proposition on dynamic algebra [8]. ■

The structure $(\text{Cong}\mathcal{M}, \wedge_{\mathcal{M}}, \vee_{\mathcal{M}})$ where, for every $\theta, \gamma \in \text{Cong}\mathcal{D}$ the operations are defined by

$$\begin{aligned}\theta \wedge_{\mathcal{M}} \gamma &= \theta \cap \gamma = (\theta_1 \cap \gamma_1, \theta_2 \cap \gamma_2) \\ \theta \vee_{\mathcal{M}} \gamma &= \langle \theta \cup \gamma \rangle_{\mathcal{M}} = (\langle \theta_1 \cup \gamma_1 \rangle_{\mathcal{B}}, \langle \theta_2 \cup \gamma_2 \rangle_{\mathcal{R}})\end{aligned}$$

is a lattice called the *congruence lattice* $\text{Cong}\mathcal{M}$ of \mathcal{M} .

In a Boolean module $\mathcal{M} = (\mathcal{B}, \mathcal{R}, :)$ we define congruences Δ_B and ∇_B on B and Δ_R and ∇_R on R as expected

$$\begin{aligned}\Delta_B &= \{(p, p) : p \in B\}, & \nabla_B &= \{(p, q) : p, q \in B\}, \\ \Delta_R &= \{(a, a) : a \in R\}, & \nabla_R &= \{(a, b) : a, b \in R\}.\end{aligned}$$

One can easily show that the pairs (Δ_B, Δ_R) , (∇_B, ∇_R) and (∇_B, Δ_R) are congruences on \mathcal{M} , but in general (Δ_B, ∇_R) is not a congruence on \mathcal{M} .

Proposition 3.3. *On a Boolean module $\mathcal{M} = (\mathcal{B}, \mathcal{R}, :)$ the pair (θ_1, ∇_R) is a modular congruence on \mathcal{M} if and only if $\theta_1 = \nabla_B$.*

Proof: If (θ_1, ∇_R) is a modular congruence on \mathcal{M} then $1\theta_1 1$ and $0\nabla_R 1$ and then $01\theta_1(1 : 1)$. So $0\theta_1 1$, i.e., $\theta_1 = \nabla_B$. \blacksquare

On an arbitrary Boolean module $\mathcal{M} = (\mathcal{B}, \mathcal{R}, :)$ (not full) it is possible that for some relation algebra elements a and b we may have $ap = bp$ for all $p \in B$ without having $a = b$. Boolean modules for which this situation is forbidden is presented next.

Definition 3.4. A Boolean module $\mathcal{M} = (\mathcal{B}, \mathcal{R}, :)$ is *bijective* if and only if, for all $a, b \in R$ we have $a = b$ whenever $ap = bp$ for all $p \in B$.

Proposition 3.5. *On a bijective Boolean module $\mathcal{M} = (\mathcal{B}, \mathcal{R}, :)$ the pair (Δ_B, θ_2) is a congruence on \mathcal{M} if and only if $\theta_2 = \Delta_R$.*

Proof: If the cardinal of the set R is 1 then Δ_R is the unique existing regular congruence. Let us admit that the cardinal of the set R is great than 1. If (Δ_B, θ_2) is a congruence and if $\theta_2 \neq \Delta_R$, then there exist distinct elements $a, b \in R$ such that $a\theta_2 b$. Immediately, $ap\Delta_B bp$ for each $p \in B$, i.e., $ap = bp$ for every $p \in B$. Since \mathcal{M} is bijective then $a = b$, a contradiction. Therefore $\theta_2 = \Delta_R$. \blacksquare

Corollary 3.6. *On a bijective Boolean module $\mathcal{M} = (\mathcal{B}, \mathcal{R}, :)$ the pair (Δ_B, ∇_R) is a congruence if and only if $\text{card } R = 1$ (if and only if $\nabla_R = \Delta_R$).*

Adopting the general classic definition of a simple algebraic structure we are able to characterize the class of simple Boolean modules.

Definition 3.7. A Boolean module $\mathcal{M} = (\mathcal{B}, \mathcal{R}, :)$ is simple whenever $\text{Cong}\mathcal{M} = \{(\Delta_B, \Delta_R), (\nabla_B, \nabla_R)\}$.

Proposition 3.8. *The degenerate Boolean module $\mathcal{M} = (\{0\}, \{o\}, :)$ is the unique simple Boolean module.*

Proof: If the cardinal of the set R is great than 1 then R admits Δ_R and ∇_R as distinct congruences. Immediately $(\Delta_B, \Delta_R), (\nabla_B, \nabla_R)$ and (∇_B, Δ_R) are different congruences on \mathcal{M} and, therefore, \mathcal{M} is not a simple Boolean module.

If the cardinal of the set R is 1, we have $R = \{o\}$ (with $o = e = 1$). Then $B = \{0\}$. In fact, by M_4 and M_5 of Definition 2.4 we have $p = ep = op = 0$ for every $p \in B$. Therefore $B = \{0\}$. \blacksquare

On a Boolean module $\mathcal{M} = (\mathcal{B}, \mathcal{R}, :)$, Boolean congruences on \mathcal{B} can exist that are not the Boolean part of any of its modular congruences. In fact, let $U = \{p, q\}$ and \mathcal{M} be the full Boolean module over U . Since $I_1 = \{\emptyset, \{p\}\}$ is a Boolean ideal on $\mathcal{B}(U)$, we can construct the Boolean congruence θ_1 on $\mathcal{B}(U)$ defined by, for $s, r \in B$, $(s, r) \in \theta_1$ if and only if $s \vee i = r \vee i$, for some $i \in I_1$, with congruence classes $[0]_{\theta_1} = \{\emptyset, \{p\}\}$ and $[q]_{\theta_1} = \{\{q\}, U\}$. Let us admit the existence of a congruence θ_2 on $\mathcal{R}(U)$ such that (θ_1, θ_2) is a modular congruence on \mathcal{M} . Let $a \in \mathcal{R}(U)$ defined by $a = \{(q, p)\}$. We have $\emptyset\theta_1\{p\}$ and $a\theta_2a$, but $(a : \emptyset, a : \{p\}) \notin \theta_1$ (in fact, $a : \emptyset = \emptyset, a : \{p\} = \{q\}$ and $(\emptyset, \{q\}) \notin \theta_1$). Therefore, on the Boolean module \mathcal{M} , $\theta_1 \in \text{Cong}\mathcal{B}$ but it does not exist a congruence $\theta_2 \in \text{Cong}\mathcal{R}$ such that $(\theta_1, \theta_2) \in \text{Cong}\mathcal{M}$.

Definition 3.9. Let $\mathcal{M} = (\mathcal{B}, \mathcal{R}, :)$ be a Boolean module. A Boolean congruence θ_1 on \mathcal{B} is called *pro-modular* on \mathcal{M} whenever there exists a congruence θ_2 on \mathcal{R} such that (θ_1, θ_2) is a modular congruence on \mathcal{M} .

Proposition 3.10. *Let $\mathcal{M} = (\mathcal{B}, \mathcal{R}, :)$ be a Boolean module and let θ_1 be a (Boolean) congruence on \mathcal{B} . The congruence θ_1 is a pro-modular congruence on \mathcal{M} if and only if the pair $(\theta_1, \Delta_{\mathcal{R}})$ is a modular congruence on \mathcal{M} .*

As previously done for dynamic algebras [7], next notion can also be established for Boolean modules.

Definition 3.11. Let $\mathcal{M} = (\mathcal{B}, \mathcal{R}, :)$ be a Boolean module with the relation algebra \mathcal{R} containing an element \exists_s satisfying $\exists_s 0 = 0$ and $\exists_s p = 1$ for every

Boolean element $p \neq 0$. This element of R is called the *simple quantifier* on \mathcal{M} .

Proposition 3.12. *If $\mathcal{M} = (\mathcal{B}, \mathcal{R}, :)$ is a Boolean module such that $\exists_s \in R$, then the congruences Δ_B and ∇_B are the only pro-modular ones.*

Proof: Let $\theta_1 \neq \Delta_B$ a pro-modular congruence on \mathcal{M} . There exists a Boolean element $p \neq 0$ such that $p\theta_1 0$. Since $\exists_s \in R$, then $\exists_s p\theta_1 \exists_s 0$, so $1\theta_1 0$. Therefore $\theta_1 = \nabla_B$. ■

Corollary 3.13. *If $\mathcal{M} = (\mathcal{B}, \mathcal{R}, :)$ is a bijective Boolean module such that $\exists_s \in R$, then the congruences (Δ_B, Δ_R) and (∇_B, θ_2) for every congruence θ_2 on \mathcal{R} , are the only modular congruences on \mathcal{M} .*

Proof: Using Propositions 3.12 and 3.5 we can infer that the congruence (Δ_B, Δ_R) is the only modular congruence with Δ_B as Boolean part. We know that, for every congruence θ_2 on \mathcal{R} , the pair (∇_B, θ_2) is a modular congruence on \mathcal{M} . ■

Corollary 3.14. *For any set U , the congruences Δ_B and ∇_B are the only pro-modular congruences in the full Boolean module over U , $\mathcal{M}(U) = (\mathcal{B}(U), \mathcal{R}(U), :)$.*

Proof: The relation ∇_R is an element of $\mathcal{R}(U)$ and ∇_R is the simple quantifier on $\mathcal{M}(U)$. ■

Proposition 3.15. *Let θ_1 be a pro-modular congruence on a Boolean module $\mathcal{M} = (\mathcal{B}, \mathcal{R}, :)$. Then*

- (1) $(\theta_1, \Delta_{\mathcal{R}})$ is the smallest modular congruence on \mathcal{M} having θ_1 as Boolean part;
- (2) $\phi = (\theta_1, \{(a, b) \in R \times R : \text{there exists } j \in R \text{ such that } a \vee j = b \vee j, j p \theta_1 0 \text{ and } j \check{p} \theta_1 0 \text{ for every } p \in B\})$ is the greatest modular congruence on \mathcal{M} having θ_1 as Boolean part.

Proof: (1) Trivial.

- (2) Our first aim is to show that $\phi = (\phi_1, \phi_2)$ with $\phi_1 = \theta_1$ and $\phi_2 = \{(a, b) \in R \times R : \text{there exists } j \in R \text{ such that } a \vee j = b \vee j, j p \theta_1 0 \text{ and } j \check{p} \theta_1 0 \text{ for every } p \in B\}$ defines a congruence on \mathcal{M} .

(a) We prove that ϕ_2 is an equivalence relation.

For $a \in R$, we have $a \vee 0 = a \vee 0$, $0 p \theta_1 0$ and $0 \check{p} \theta_1 0$, and therefore $a \phi_2 a$.

Trivially if $a\phi_2b$ then $b\phi_2a$.

Let $a\phi_2b$ and $b\phi_2c$. So there exist $j, k \in R$ such that $a \vee j = b \vee j$, $jp\theta_1 0$ and $j\check{p}\theta_1 0$ for every $p \in B$ and $b \vee k = c \vee k$, $kp\theta_1 0$ and $k\check{p}\theta_1 0$ for every $p \in B$.

But $a \vee (j \vee k) = (a \vee j) \vee k = (b \vee j) \vee k = (b \vee k) \vee j = (c \vee k) \vee j = c \vee (j \vee k)$.

Since $jp\theta_1 0$ and $kp\theta_1 0$ then $jp \vee kp\theta_1 0$ so $(j \vee k)p\theta_1 0$.

Since $j\check{p}\theta_1 0$ and $k\check{p}\theta_1 0$ then $j\check{p} \vee k\check{p}\theta_1 0$ so $(j\check{\vee} k\check{)}p\theta_1 0$, i.e., $(j \vee k)\check{p}\theta_1 0$.

Therefore $a\phi_2c$.

- (b) Let $a, b, c, d \in R$ such that $a\phi_2b$ and $c\phi_2d$. We have to prove that $a\check{\phi}_2b\check{}$, $(a \wedge c)\phi_2(b \wedge d)$, $(a \vee c)\phi_2(b \vee d)$ and $ac\phi_2bd$.

Since $a\phi_2b$ and $c\phi_2d$ there exist $j, k \in R$ such that $a \vee j = b \vee j$, $jp\theta_1 0$ and $j\check{p}\theta_1 0$ for every $p \in B$ and $c \vee k = d \vee k$, $kp\theta_1 0$ and $k\check{p}\theta_1 0$ for every $p \in B$.

(i) We have $a\check{\vee} j\check{=} = (a \vee j)\check{=} = (b \vee j)\check{=} = b\check{\vee} j\check{}$ and $j\check{p}\theta_1 0$ and $j\check{\check{p}} = jp\theta_1 0$ so $a\check{\phi}_2b\check{}$.

(ii) We have $(a \wedge c) \vee [(j \wedge c) \vee (a \wedge k) \vee (j \wedge k)] = (a \vee j) \wedge (c \vee k) = (b \vee j) \wedge (d \vee k) = (b \wedge d) \vee [(j \wedge d) \vee (b \wedge k) \vee (j \wedge k)]$. Let $m = (j \wedge c) \vee (a \wedge k) \vee (j \wedge k)$ and $n = (j \wedge d) \vee (b \wedge k) \vee (j \wedge k)$. So $m\check{=} = (j\check{\wedge} c\check{)} \vee (a\check{\wedge} k\check{)} \vee (j\check{\wedge} k\check{)}$ and $n\check{=} = (j\check{\wedge} d\check{)} \vee (b\check{\wedge} k\check{)} \vee (j\check{\wedge} k\check{)}$. Since $(a \wedge c) \vee m = (b \wedge d) \vee n$ then $(a \wedge c) \vee (m \vee n) = (b \wedge d) \vee (m \vee n)$.

Since for every $p \in B$, $(j \wedge c)p \leq jp$, $(a \wedge k)p \leq kp$, $(j \wedge k)p \leq jp$, $(j \wedge d)p \leq jp$, $(b \wedge k)p \leq kp$, $jp\theta_1 0$ and $kp\theta_1 0$ then $(j \wedge c)p\theta_1 0$, $(a \wedge k)p\theta_1 0$, $(j \wedge k)p\theta_1 0$, $(j \wedge d)p\theta_1 0$, $(b \wedge k)p\theta_1 0$. So $mp\theta_1 0$ and $np\theta_1 0$, and therefore $(m \vee n)p\theta_1 0$.

Since for every $p \in B$, $(j\check{\wedge} c\check{)}p \leq j\check{p}$, $(a\check{\wedge} k\check{)}p \leq k\check{p}$, $(j\check{\wedge} k\check{)}p \leq j\check{p}$, $(j\check{\wedge} d\check{)}p \leq j\check{p}$, $(b\check{\wedge} k\check{)}p \leq k\check{p}$, $j\check{p}\theta_1 0$ and $k\check{p}\theta_1 0$ then $(j\check{\wedge} c\check{)}p\theta_1 0$, $(a\check{\wedge} k\check{)}p\theta_1 0$, $(j\check{\wedge} k\check{)}p\theta_1 0$, $(j\check{\wedge} d\check{)}p\theta_1 0$, $(b\check{\wedge} k\check{)}p\theta_1 0$. So $m\check{p}\theta_1 0$ and $n\check{p}\theta_1 0$, and so $(m \vee n)\check{p} = (m\check{\vee} n\check{)}p\theta_1 0$.

Therefore $(a \wedge c)\phi_2(b \wedge d)$.

- (iii) We have $(a \vee c) \vee (j \vee k) = (b \vee d) \vee (j \vee k)$.

Since $jp\theta_1 0$, $kp\theta_1 0$ then $jp \vee kp\theta_1 0$, i.e., $(j \vee k)p\theta_1 0$.

Since $j\check{p}\theta_1 0$, $k\check{p}\theta_1 0$ then $j\check{p} \vee k\check{p}\theta_1 0$, i.e., $(j\check{\vee} k\check{)}p\theta_1 0$. So $(j \vee k)\check{p}\theta_1 0$. Therefore $(a \vee c)\phi_2(b \vee d)$.

(iv) We have $ac \vee (ak \vee jd) = (ac \vee ak) \vee jd = a(c \vee k) \vee jd = a(d \vee k) \vee jd = (ad \vee ak) \vee jd = (a \vee j)d \vee ak = (b \vee j)d \vee ak = bd \vee (jd \vee ak)$. Let $l = ak \vee jd$. So $l^\sim = (ak)^\sim \vee (jd)^\sim = k^\sim a^\sim \vee d^\sim j^\sim$. Since $kp\theta_1 0$ and $jp\theta_1 0$ for every $p \in B$ then $akp\theta_1 0$ and $jd\check{p}\theta_1 0$ for every $p \in B$. So $(akp \vee jd\check{p})\theta_1 0$, i.e., $(ak \vee jd)p\theta_1 0$, i.e., $lp\theta_1 0$. Since $k^\sim p\theta_1 0$ and $j^\sim p\theta_1 0$ for every $p \in B$ then $k^\sim a^\sim p\theta_1 0$ and $d^\sim j^\sim p\theta_1 0$ for every $p \in B$. So $(k^\sim a^\sim p \vee d^\sim j^\sim p)\theta_1 0$, i.e., $(k^\sim a^\sim \vee d^\sim j^\sim)p\theta_1 0$, i.e., $l^\sim p\theta_1 0$.

Therefore $ac\phi_2 bd$.

(c) Let $p\theta_1 q$ and $a\phi_2 b$. We have to prove that $ap\theta_1 bq$.

Since $a\phi_2 b$ then there exists $j \in R$ such that $a \vee j = b \vee j$, $jp\theta_1 0$ and $j^\sim p\theta_1 0$ for every $p \in B$.

So $(a \vee j)p = (b \vee j)p$, i.e., $ap \vee jp = bp \vee jp$. Since $jp\theta_1 0$ then $(ap \vee jp)\theta_1 (0 \vee ap)$ and $(bp \vee jp)\theta_1 (0 \vee bp)$, i.e., $(ap \vee jp)\theta_1 ap$ and $(bp \vee jp)\theta_1 bp$. So $ap\theta_1 bp$. Since $p\theta_1 q$ and θ_1 is pro-modular then $bp\theta_1 bq$. Therefore $ap\theta_1 bq$.

(d) Now we have to prove that ϕ is the smallest modular congruence having θ_1 as Boolean part, i.e., if $\theta = (\theta_1, \theta_2)$ is a modular congruence on \mathcal{M} , then $\theta \subseteq \phi$.

Let $a, b \in R$ and $a\theta_2 b$. Since θ is a modular congruence on \mathcal{M} we have $a'\theta_2 b'$ and $a^\sim\theta_2 b^\sim$. So $(a' \wedge b)\theta_2 0$ and $(a \wedge b')\theta_2 0$ and then $[(a' \wedge b) \vee (a \wedge b')]\theta_2 0$. Therefore $(a \oplus b)\theta_2 0$ and $(a \oplus b)p\theta_1 0$.

We also have $(a' \wedge b^\sim)\theta_2 0$ and $(a^\sim \wedge b')\theta_2 0$ and then $[(a' \wedge b^\sim) \vee (a^\sim \wedge b')]\theta_2 0$. Therefore $(a^\sim \oplus b^\sim)\theta_2 0$, i.e., $(a \oplus b)\theta_2 0$ and $(a \oplus b)^\sim p\theta_1 0$. Since $a \vee (a \oplus b) = b \vee (a \oplus b)$ then $a\phi_2 b$. ■

Definition 3.16. Let θ_1 be a pro-modular congruence on a Boolean module $\mathcal{M} = (\mathcal{B}, \mathcal{R}, :)$ and let $\phi = (\phi_1, \phi_2)$ be defined by

$$\phi_1 = \theta_1$$

$$\phi_2 = \{(a, b) \in R \times R : \text{there exists } j \in R \text{ such that } a \vee j = b \vee j, jp\theta_1 0 \text{ and } j^\sim p\theta_1 0 \text{ for every } p \in B\}.$$

The relation ϕ is called the *determining congruence* of any $\theta \in \text{Cong}\mathcal{D}$ having θ_1 as Boolean part (or simply a *determining congruence*).

Next example illustrates Proposition 3.15.

Example 3.17. Let $\mathcal{M} = (\mathcal{B}, \mathcal{R}, :)$ be the proper Boolean module where $B = \{\emptyset, \{p\}, \{q\}, \{p, q\}\}$, $R = \{\Lambda, a, b, c\}$, Λ is the empty relation, $a = \{(p, p)\}$,

$b = \{(q, q)\}$ and $c = \{(p, p), (q, q)\}$. Let θ_1 be the Boolean congruence with congruence classes $[\emptyset]_{\theta_1} = \{\emptyset, \{p\}\}$ and $[\{p, q\}]_{\theta_1} = \{\{q\}, \{p, q\}\}$, i.e.,

$\theta_1 = \{(\emptyset, \emptyset), (\{p\}, \{p\}), (\{q\}, \{q\}), (\{p, q\}, \{p, q\}), (\emptyset, \{p\}), (\{p\}, \emptyset), (\{q\}, \{p, q\}), (\{p, q\}, \{q\})\}$.

We have

$$\begin{array}{ll} \Lambda \check{\emptyset} = \Lambda \emptyset = \emptyset & a \check{\emptyset} = a \emptyset = \emptyset \\ \Lambda \check{\{p\}} = \Lambda \{p\} = \emptyset & a \check{\{p\}} = a \{p\} = \{p\} \\ \Lambda \check{\{q\}} = \Lambda \{q\} = \emptyset & a \check{\{q\}} = a \{q\} = \emptyset \\ \Lambda \check{\{p, q\}} = \Lambda \{p, q\} = \emptyset & a \check{\{p, q\}} = a \{p, q\} = \{p\} \end{array}$$

$$\begin{array}{ll} b \check{\emptyset} = b \emptyset = \emptyset & c \check{\emptyset} = c \emptyset = \emptyset \\ b \check{\{p\}} = b \{p\} = \emptyset & c \check{\{p\}} = c \{p\} = \{p\} \\ b \check{\{q\}} = b \{q\} = \{q\} & c \check{\{q\}} = c \{q\} = \{q\} \\ b \check{\{p, q\}} = b \{p, q\} = \{q\} & c \check{\{p, q\}} = c \{p, q\} = \{p, q\} \end{array}$$

So (θ_1, Δ_R) is a modular congruence on \mathcal{M} and then is the smallest (modular) congruence on \mathcal{M} having θ_1 as Boolean part.

The greatest modular congruence on \mathcal{M} having θ_1 as Boolean part is (θ_1, ϕ_2) for $\phi_2 = \{(f, g) \in R \times R : \text{there exists } j \in R \text{ such that } f \vee j = g \vee j, js\theta_1\emptyset \text{ and } j\check{s}\theta_1\emptyset \text{ for every } s \in B\}$.

So Λ and a are the only elements j of R such that $js\theta_1\emptyset$ and $j\check{s}\theta_1\emptyset$ for every $s \in B$. Trivially we have $s \vee \Lambda = s \vee \Lambda$ for every $s \in R$, $\Lambda \vee a = a \vee a$, $b \vee a = c \vee a$ and for every $j \in \{\Lambda, a\}$ we have $\Lambda \vee j \neq b \vee j$, $\Lambda \vee j \neq c \vee j$, $a \vee j \neq b \vee j$ and $a \vee j \neq c \vee j$. So $\phi_2 = \{(\Lambda, \Lambda), (a, a), (b, b), (c, c), (\Lambda, a), (a, \Lambda), (b, c), (c, b)\}$.

4. The lattice $Ide\mathcal{M}$ of modular ideals

Usually, the notion of *ideal* in a given class of algebras is established so that the zero-classes of congruence relations are easily seen to be ideals.

Definition 4.1. A (*modular*) *ideal* on a Boolean module $\mathcal{M} = (\mathcal{B}, \mathcal{R}, :)$ is a pair $I = (I_1, I_2)$ satisfying the following conditions

- (1) I_1 is a Boolean ideal on \mathcal{B} ;
- (2) If $p \in I_1$ and $a \in R$ then $ap \in I_1$;
- (3) (a) I_2 is a Boolean ideal on \mathcal{R} ;
(b) If $a \in I_2, c \in R$ then $ac, ca, a \check{\vee} \in I_2$;
- (4) If $a \in I_2$ and $p \in B$, then $ap \in I_1$.

Such a subset I_2 of R satisfying condition (3) is called an *ideal* of R .

We denote by $Ide\mathcal{M}$ the set of all ideals on a Boolean module $\mathcal{M} = (\mathcal{B}, \mathcal{R}, :)$. We intend to insert a lattice structure into $Ide\mathcal{M}$. To do so we need to define, for arbitrary modular ideals I and J , $I \wedge_{\mathcal{M}} J$ and $I \vee_{\mathcal{M}} J$. It is immediate to put $I \wedge_{\mathcal{M}} J = (I_1 \cap J_1, I_2 \cap J_2)$. Once again the disjunction requires some attention. We denote by $\langle X \rangle_{\mathcal{A}}$ the ideal generated by a subset X of any (homogeneous or heterogeneous) algebra \mathcal{A} , i.e., the intersection of all ideals I on \mathcal{A} containing X ,

$$\langle X \rangle_{\mathcal{A}} = \cap \{I : I \text{ ideal on } \mathcal{A} \text{ and } X \subseteq I\}.$$

Proposition 4.2. *Let $\mathcal{M} = (\mathcal{B}, \mathcal{R}, :)$ a Boolean module, $I = (I_1, I_2)$ and $J = (J_1, J_2)$ be elements of $Ide\mathcal{M}$. We have*

$$\begin{aligned} \langle I_1 \cup J_1 \rangle_{\mathcal{B}} &= \{p \in B : p \leq p_1 \vee p_2, \text{ for some } p_i \in I_1 \cup J_1, i = 1, 2\} \\ \langle I_2 \cup J_2 \rangle_{\mathcal{R}} &= \{a \in R : a \leq a_1 \vee a_2, \text{ for some } a_i \in I_2 \cup J_2, i = 1, 2\}. \end{aligned}$$

Proof: We only have to prove the second identity since the first is a well known Boolean algebras result [4]. We want to show that, for $X = \{a \in R : a \leq a_1 \vee a_2, \text{ for some } a_i \in I_2 \cup J_2, i = 1, 2\}$, we have

- (i) X is an ideal of \mathcal{R} ;
- (ii) $I_2 \cup J_2 \subseteq X$;
- (iii) if Y is an ideal of \mathcal{R} and $I_2 \cup J_2 \subseteq Y$, then $X \subseteq Y$.

A well known Boolean algebras result states that X is a Boolean ideal on R . Now, let $a \in X, b \in R$. Then $a \leq a_1 \vee a_2$, for some $a_i \in I_2 \cup J_2, i = 1, 2$. So $ca \leq ca_1 \vee ca_2$, $ac \leq a_1c \vee a_2c$ and $a^\smile \leq a_1^\smile \vee a_2^\smile$. Since for $i = 1, 2$, $ca_i, a_i c$ and a_i^\smile are elements of $I_2 \cup J_2$, we get ca, ac and a^\smile in $\langle I_2 \cup J_2 \rangle_{\mathcal{R}}$. Therefore X is an ideal of \mathcal{R} . It is straightforward that $I_2 \cup J_2 \subseteq X$. Let $a \in X$ and Y be an ideal of \mathcal{R} such that $I_2 \cup J_2 \subseteq Y$. Then $a \leq a_1 \vee a_2$, for some $a_i \in I_2 \cup J_2 \subseteq Y, i = 1, 2$. But Y is an ideal of \mathcal{R} and $a_i \in Y$ for $i = 1, 2$ so $a_1 \vee a_2 \in Y$. Therefore $a \leq a_1 \vee a_2 \in Y$. Since Y is an ideal of \mathcal{R} we get $a \in Y$. ■

For $\mathcal{M} = (\mathcal{B}, \mathcal{R}, :)$ Boolean module, $I \subseteq B$ and $J \subseteq R$ we write JI to represent the set $JI = \{ap : a \in J \text{ and } p \in I\}$.

Proposition 4.3. *Let $\mathcal{M} = (\mathcal{B}, \mathcal{R}, :)$ be a Boolean module, $I = (I_1, I_2)$ and $J = (J_1, J_2)$ elements of $Ide\mathcal{M}$. We have*

$$\begin{aligned} R(\langle I_1 \cup J_1 \rangle_{\mathcal{B}}) &\subseteq \langle I_1 \cup J_1 \rangle_{\mathcal{B}}, \\ (\langle I_2 \cup J_2 \rangle_{\mathcal{R}})B &\subseteq \langle I_1 \cup J_1 \rangle_{\mathcal{B}}. \end{aligned}$$

Proof: Analogous to proposition on dynamic algebra [8]. ■

Therefore the structure $\mathcal{I}de\mathcal{M} = (Ide\mathcal{M}, \wedge_{\mathcal{M}}, \vee_{\mathcal{M}})$ where, for every $I = (I_1, I_2), J = (J_1, J_2) \in Ide\mathcal{M}$, the operations are defined by

$$\begin{aligned} I \wedge_{\mathcal{M}} J &= I \cap J = (I_1 \cap J_1, I_2 \cap J_2) \\ I \vee_{\mathcal{M}} J &= \langle I \cup J \rangle_{\mathcal{M}} = (\langle I_1 \cup J_1 \rangle_{\mathcal{B}}, \langle I_2 \cup J_2 \rangle_{\mathcal{R}}) \end{aligned}$$

is a lattice, called the *lattice of ideals* of \mathcal{M} .

Similarly to the congruences case, on a Boolean module $\mathcal{M} = (\mathcal{B}, \mathcal{R}, :)$, Boolean ideals on \mathcal{B} can exist that are not the Boolean part of any modular ideal on \mathcal{M} . In fact, let $U = \{p, q\}$ and \mathcal{M} the full Boolean module over U . The set $I_1 = \{\emptyset, \{p\}\}$ is a Boolean ideal on $\mathcal{B}(U)$ but, since for $a \in \mathcal{R}(U)$ given by $a = \{(q, p)\}$ we have $a : \{p\} = \{q\} \notin I_1$, the pair (I_1, I_2) is not a modular ideal on \mathcal{M} , for any subset I_2 of \mathcal{R} (by 2 of Definition 4.1). Thus we are led to establish the following definition.

Definition 4.4. Let $\mathcal{M} = (\mathcal{B}, \mathcal{R}, :)$ be a Boolean module. A Boolean ideal I_1 on \mathcal{B} is called *pro-modular* on \mathcal{M} if there exists an ideal I_2 of \mathcal{R} such that (I_1, I_2) is a modular ideal on \mathcal{M} .

Proposition 4.5. Let $\mathcal{M} = (\mathcal{B}, \mathcal{R}, :)$ be a Boolean module and let I_1 be a Boolean ideal on \mathcal{B} . The ideal I_1 is a pro-modular ideal on \mathcal{M} if and only if $(I_1, \{o\})$ is a modular ideal on \mathcal{M} .

Proposition 4.6. Let I_1 be a pro-modular ideal on a Boolean module $\mathcal{M} = (\mathcal{B}, \mathcal{R}, :)$. Then

- (1) $(I_1, \{o\})$ is the smallest modular ideal having I_1 as Boolean part;
- (2) $F = (I_1, \{a : ap, a\check{p} \in I_1 \text{ for every } p \in B\})$ is the greatest modular ideal having I_1 as Boolean part.

Proof: (1) It is trivial that $(I_1, \{o\})$ is the smallest modular ideal having I_1 as Boolean part.

- (2) Let $F = (I_1, F_2)$ with $F_2 = \{a : ap, a\check{p} \in I_1 \text{ for every } p \in B\}$.

(a) I_1 is a Boolean ideal on \mathcal{B} ;

(b) Since I_1 is pro-modular ideal, for $p \in I_1$ and $a \in R$ we have $ap \in I_1$;

(c) (i) Since $o \in F_2$ ($o\check{p} = op = 0$ for every $p \in B$), $F_2 \neq \emptyset$.

Let $a, b \in F_2$ and $d \in R$ such that $d \leq a$. So $ap, a\check{p}, bp, b\check{p} \in I_1$ for every $p \in B$. But

$(a \vee b)p = ap \vee bp \in I_1$ and $(a \vee b)\check{p} = (a\check{p} \vee b\check{p}) \in I_1$, so $a \vee b \in F_2$.

For every $p \in B$ we have $dp \leq ap \in I_1$, so $dp \in I_1$ and $d\check{p} \leq a\check{p} \in I_1$, so $d\check{p} \in I_1$, and therefore $d \in F_2$.

So F_2 is a Boolean ideal on \mathcal{R} ;

(ii) Let $a \in F_2$ and $c \in R$. So $ap, a\check{p} \in I_1$ for every $p \in B$.

Then

$(ac)p = a(cp) \in I_1$ and $(ac)\check{p} = (c\check{a})p = c\check{(a\check{p})} \in I_1$, so $ac \in F_2$.

$(ca)p = c(ap) \in I_1$ and $(ca)\check{p} = (a\check{c})p = a\check{(c\check{p})} \in I_1$, so $ca \in F_2$.

$a\check{p} \in I_1$ and $a\check{\check{p}} = ap \in I_1$, so $a\check{\check{p}} \in F_2$.

(d) By definition of F_2 , if $a \in F_2$ and $p \in B$, then $ap \in I_1$. Therefore, (I_1, F_2) is a modular ideal on \mathcal{M} .

Let $I = (I_1, I_2)$ be an arbitrary modular ideal on \mathcal{M} and $a \in I_2$. Then $a\check{\check{p}} \in I_2$ and $ap \in I_1$ for every $p \in B$. We also have $a\check{p} \in I_1$ for every $p \in B$ establishing the conditions required to $a \in F_2$. Therefore, (I_1, F_2) is the greatest modular ideal on \mathcal{M} having I_1 as Boolean part. ■

Definition 4.7. Let I_1 be a pro-modular ideal on a Boolean module $\mathcal{M} = (\mathcal{B}, \mathcal{R}, :)$ and let $F = (F_1, F_2)$ be defined by

$$F_1 = I_1$$

$$F_2 = \{a : ap, a\check{p} \in I_1 \text{ for every } p \in B\}.$$

We say that F is the *determining ideal* of any $I \in Ide\mathcal{M}$ having I_1 as Boolean part (or simply, a *determining ideal*).

Next example illustrates Proposition 4.6.

Example 4.8. Let $\mathcal{M} = (\mathcal{B}, \mathcal{R}, :)$ be the Boolean module defined in Example 3.17, i.e., $B = \{\emptyset, \{p\}, \{q\}, \{p, q\}\}$, $R = \{\Lambda, a, b, c\}$, Λ is the empty relation, $a = \{(p, p)\}$, $b = \{(q, q)\}$ and $c = \{(p, p), (q, q)\}$. Let I_1 be the Boolean ideal $I_1 = \{\emptyset, \{p\}\}$.

We have

$$\begin{array}{ll} \Lambda\check{\emptyset} = \Lambda\emptyset = \emptyset & a\check{\emptyset} = a\emptyset = \emptyset \\ \Lambda\check{\{p\}} = \Lambda\{p\} = \emptyset & a\check{\{p\}} = a\{p\} = \{p\} \\ \Lambda\check{\{q\}} = \Lambda\{q\} = \emptyset & a\check{\{q\}} = a\{q\} = \emptyset \\ \Lambda\check{\{p, q\}} = \Lambda\{p, q\} = \emptyset & a\check{\{p, q\}} = a\{p, q\} = \{p\} \end{array}$$

$$\begin{aligned}
 b\check{\emptyset} &= b\emptyset = \emptyset & c\check{\emptyset} &= c\emptyset = \emptyset \\
 b\check{\{p\}} &= b\{p\} = \emptyset & c\check{\{p\}} &= c\{p\} = \{p\} \\
 b\check{\{q\}} &= b\{q\} = \{q\} & c\check{\{q\}} &= c\{q\} = \{q\} \\
 b\check{\{p, q\}} &= b\{p, q\} = \{q\} & c\check{\{p, q\}} &= c\{p, q\} = \{p, q\}
 \end{aligned}$$

So $(I_1, \{\Lambda\})$ is a modular ideal on \mathcal{M} and thus is the smallest ideal on \mathcal{M} having I_1 as Boolean part.

The greatest modular ideal on \mathcal{M} having I_1 as Boolean part is (I_1, F_2) with $F_2 = \{f \in R : fs, f\check{s} \in I_1 \text{ for every } s \in B\}$ and Λ and a are the only elements j of R such that $js, j\check{s} \in I_1$ for every $s \in B$. Therefore $F_2 = \{\Lambda, a\}$.

5. Modular Congruences and Modular Ideals

The main purpose of this paragraph is to establish that the class of Boolean module is ideal determined [5], i.e., that each modular ideal is the zero-class of a unique modular congruence.

Definition 5.1. If $\theta = (\theta_1, \theta_2) \in \text{Cong}\mathcal{M}$ where $\mathcal{M} = (\mathcal{B}, \mathcal{R}, :)$ is a Boolean module, we say that $\mathcal{I}(\theta) = \mathcal{I}^\theta = (\mathcal{I}_1^\theta, \mathcal{I}_2^\theta)$ defined by

$$\begin{aligned}
 \mathcal{I}_1^\theta &= \{p \in B : p\theta_1 0\} = [0]_{\theta_1} \\
 \mathcal{I}_2^\theta &= \{a \in R : a\theta_2 0\} = [0]_{\theta_2}
 \end{aligned}$$

is the *kernel of the congruence* θ .

Proposition 5.2. *The kernel $\mathcal{I}(\theta)$ of a congruence θ on a Boolean module $\mathcal{M} = (\mathcal{B}, \mathcal{R}, :)$ is an ideal on \mathcal{M} .*

Proof: (1) The fact that \mathcal{I}_1^θ is a Boolean ideal on \mathcal{B} is a known Boolean algebras result.

(2) We have to prove that if $p \in \mathcal{I}_1^\theta$ and $a \in R$ then $ap \in \mathcal{I}_1^\theta$. In fact, if $a \in R$ and $p \in \mathcal{I}_1^\theta$, then $p\theta_1 0$ and $a\theta_2 a$. Therefore $(ap)\theta_1(a0)$, i.e., $(ap)\theta_1 0$. So $ap \in \mathcal{I}_1^\theta$.

(3) (a) The fact that \mathcal{I}_2^θ is a Boolean ideal on \mathcal{R} is again a known Boolean algebras result.

(b) We have to prove that, if $a \in \mathcal{I}_2^\theta$ and $c \in R$, then $ac, ca, a\check{\in} \in \mathcal{I}_2^\theta$. In fact, since $a \in \mathcal{I}_2^\theta$ then $a\theta_2 0$, and therefore, since $c\theta_2 c$ then $(ac)\theta_2(oc)$ and $(ca)\theta_2(co)$, i.e., $(ac)\theta_2 0$ and $(ca)\theta_2 0$. To this extend ac and ca are elements of \mathcal{I}_2^θ . Since $a\theta_2 0$ and θ_2 is a congruence on \mathcal{R} then $a\check{\theta}_2 0\check{\in}$, so $a\check{\in}\theta_2 0$, and therefore $a\check{\in} \in \mathcal{I}_2^\theta$.

(4) Let $a \in \mathcal{I}_2^\theta$ and $p \in B$. Then $a\theta_2 0$ and $p\theta_1 p$, and therefore $(ap)\theta_1 op$, i.e., $(ap)\theta_1 0$. Immediately $ap \in \mathcal{I}_1^\theta$. ■

Definition 5.3. The *kernel* of a Boolean modular homomorphism $h = (h_1, h_2) : \mathcal{M} \longrightarrow \mathcal{M}'$ between Boolean modules is the pair $(\{p \in B : h_1(p) = 0\}, \{a \in R : h_2(a) = 0\})$.

Proposition 5.4. *The kernel of a Boolean modular homomorphism $h : \mathcal{M} \longrightarrow \mathcal{M}'$ between Boolean modules is a modular ideal on \mathcal{M} .*

Proof: Trivial. ■

Definition 5.5. If $I = (I_1, I_2)$ is a modular ideal on a Boolean module $\mathcal{M} = (\mathcal{B}, \mathcal{R}, :)$, we define $\mathcal{C}(I) = \mathcal{C}^I = (\mathcal{C}_1^I, \mathcal{C}_2^I)$ by

$$\begin{aligned} p \mathcal{C}_1^I q &\text{ if and only if } p \vee i = q \vee i \text{ for some } i \in I_1, \\ a \mathcal{C}_2^I b &\text{ if and only if } a \vee j = b \vee j \text{ for some } j \in I_2, \end{aligned}$$

for $p, q \in B$ and $a, b \in R$.

Proposition 5.6. *If $I = (I_1, I_2)$ is a modular ideal on a Boolean module $\mathcal{M} = (\mathcal{B}, \mathcal{R}, :)$, then $\mathcal{C}(I)$ is a modular congruence on \mathcal{M} .*

Proof: (i) The relation \mathcal{C}_1^I is a congruence relation on \mathcal{B} , a known result in Boolean algebras.

(ii) The relation \mathcal{C}_2^I is a Boolean congruence relation on \mathcal{R} , a known result in Boolean algebras. To prove that the relation \mathcal{C}_2^I is a congruence relation on \mathcal{R} we have to prove that, for $a, b, c, d \in R$, if $a \mathcal{C}_2^I b$ and $c \mathcal{C}_2^I d$ then $(ac) \mathcal{C}_2^I (bd)$ and $(a \smile) \mathcal{C}_2^I (b \smile)$. Let us admit that $a \mathcal{C}_2^I b$ and $c \mathcal{C}_2^I d$. Then there exist j, k in I_2 such that $a \vee j = b \vee j$ and $c \vee k = d \vee k$. Now

(a) From $c \vee k = d \vee k$ we get $a(c \vee k) = a(d \vee k)$, i.e., $ac \vee ak = ad \vee ak$. Hence $ac \vee (ak \vee jd) = a(c \vee k) \vee jd = a(d \vee k) \vee jd = (ad \vee ak) \vee jd = (a \vee j)d \vee ak = (b \vee j)d \vee ak = bd \vee (jd \vee ak)$. Since $jd \vee ak \in I_2$ we get $(ac) \mathcal{C}_2^I (bd)$.

(b) We have $a \smile \vee j \smile = (a \vee j) \smile = (b \vee j) \smile = b \smile \vee j \smile$. Since $j \smile \in I_2$ then $(a \smile) \mathcal{C}_2^I (b \smile)$.

(iii) Now we have to prove that, for $a, b \in R$ and $p, q \in B$, if $a \mathcal{C}_2^I b$ and $p \mathcal{C}_1^I q$ then $(ap) \mathcal{C}_1^I (bq)$. Since $p \mathcal{C}_1^I q$ and $a \mathcal{C}_2^I b$, then $p \vee i = q \vee i$ for some $i \in I_1$ and $a \vee j = b \vee j$ for some $j \in I_2$ (and therefore $aq \vee jq = bq \vee jq$). But from $p \vee i = q \vee i$ we get $ap \vee ai = aq \vee ai$ and moreover $ap \vee ai \vee jq = aq \vee ai \vee jq$. So $ap \vee (ai \vee jq) = (ap \vee ai) \vee jq = (aq \vee ai) \vee jq = (aq \vee jq) \vee ai = (bq \vee jq) \vee ai = bq \vee (ai \vee jq)$. Since $ai \vee jq \in I_1$ then $(ap) \mathcal{C}_1^I (bq)$. ■

Definition 5.7. For I a modular ideal on a Boolean module, we say that $\mathcal{C}(I)$ is the *congruence induced* by I .

Proposition 5.8. *If $I = (I_1, I_2)$ is a modular ideal on a Boolean module, then $\mathcal{I}(\mathcal{C}(I)) = I$.*

Proof: Similar to Boolean algebras. ■

Proposition 5.9. *On a Boolean module a modular ideal is a determining ideal if and only if it is the kernel of a determining congruence.*

Proof: Let $I = (I_1, I_2)$ be a modular ideal on a Boolean module $\mathcal{M} = (\mathcal{B}, \mathcal{R}, :)$. By Prop.5.8 there exists a modular congruence θ such that $I = \mathcal{I}(\theta)$. Let ϕ be the determining congruence of θ . We have $\phi_1 = \theta_1$ and $\phi_2 = \{(a, b) \in R \times R : \text{there exists } j \in R \text{ such that } a \vee j = b \vee j, jp\theta_1 0 \text{ and } j\check{p}\theta_1 0 \text{ for every } p \in B\}$. So, $\mathcal{I}_1^\phi = \mathcal{I}_1^\theta = I_1$ and $\mathcal{I}_2^\phi = \{a : a\phi_2 0\}$. Then, we have to prove that $\{a : a\phi_2 0\} = \{a : ap, a\check{p} \in I_1 \text{ for every } p \in B\}$.

Let $a_1 \in \{a : a\phi_2 0\}$. There exists $j \in R$ such that $a_1 \vee j = 0 \vee j, jp\theta_1 0$ and $j\check{p}\theta_1 0$ for every $p \in B$. Since $a_1 \vee j = 0 \vee j = j$ then $a_1 p \vee jp = jp$ and $a_1 \check{p} \vee j\check{p} = j\check{p}$ so $a_1 \check{p} \vee j\check{p} = j\check{p}$. But $a_1 p \leq a_1 p \vee jp = jp\theta_1 0$ so $a_1 p\theta_1 0$. Similarly, we have $a_1 \check{p} \leq a_1 \check{p} \vee j\check{p} = j\check{p}\theta_1 0$ so $a_1 \check{p}\theta_1 0$ and then $a_1 p, a_1 \check{p} \in [0]_{\theta_1} = I_1$.

Let $a_1 \in \{a : ap, a\check{p} \in I_1 \text{ for every } p \in B\}$. We have to prove that $a_1\phi_2 0$, i.e., there exists $j \in R$ such that $a_1 \vee j = j, jp\theta_1 0$ and $j\check{p}\theta_1 0$ for every $p \in B$. Since $a_1 p, a_1 \check{p} \in I_1$ for every $p \in B$ and $I_1 = [0]_{\theta_1}$ then $a_1 p\theta_1 0$ and $a_1 \check{p}\theta_1 0$. Since $a_1 \vee a_1 = a_1$ putting $j = a_1$ we have the required. ■

Proposition 5.10. *If $\theta = (\theta_1, \theta_2)$ is a congruence on a Boolean module, then $\mathcal{C}(\mathcal{I}(\theta)) = \theta$.*

Proof: As in Boolean algebras. ■

Proposition 5.11. *On a Boolean module a modular congruence is a determining congruence if and only if it is the congruence induced by a determining ideal.*

Proof: If ϕ is a determining congruence on a Boolean module, Proposition 5.9 asserts that $\mathcal{I}(\phi) = F$ for some determining ideal F . So $\mathcal{C}(\mathcal{I}(\phi)) = \mathcal{C}(F)$. But Proposition 5.10 infers that $\phi = \mathcal{C}(F)$.

If F is a determining ideal on a Boolean module, using Proposition 5.9 we have $\mathcal{I}(\phi) = F$ for some determining congruence ϕ . So $\mathcal{C}(\mathcal{I}(\phi)) = \mathcal{C}(F)$. By Proposition 5.10 we have $\phi = \mathcal{C}(F)$ as required. ■

Theorem 5.12. *The pair of maps $\mathcal{C} : Ide\mathcal{M} \rightarrow Cong\mathcal{M}$ (that for each $I \in Ide\mathcal{M}$ assigns the congruence $\mathcal{C}(I)$) and $\mathcal{I} : Cong\mathcal{M} \rightarrow Ide\mathcal{M}$ (that for each $\theta \in Cong\mathcal{M}$ assigns the ideal $\mathcal{I}(\theta)$) defines an isomorphism between the lattices $Ide\mathcal{M}$ and $Cong\mathcal{M}$.*

Proof: As in Boolean algebras. ■

We infer that the class of the Boolean module is ideal determined, i.e., each ideal is the zero-class of a unique congruence. We can easily affirm that the modular ideal F defined on Proposition 4.6 is the kernel of the modular congruence ϕ presented on Proposition 3.15. Conversely, the congruence ϕ defined on Proposition 3.15 is the congruence induced by the modular ideal F constructed on Proposition 4.6, i.e., either Proposition 4.6 and Proposition 3.15 can now be stated as corollaries of each other using Theorem 5.12.

References

- [1] C. Brink. *Boolean Modules*, Journal of Algebra **71**, 291-313 (1981).
- [2] C. Brink, K. Britz, and R. Schmidt. *Peirce Algebras*, Formal Aspects of Computing **6**, 339-358 (1994).
- [3] L. H. Chin, and A. Tarski. *Distributive and modular laws in the arithmetic of relation algebras*, University of California Publications in Mathematics **1**, 341-384 (1951).
- [4] G. Gratzer. *General Lattice Theory*, (Birkhauser Verlag, 1998).
- [5] H.P. Gumm, and A. Ursini. *Ideals in universal algebras*, Algebra Universalis **19**, 45-54 (1984).
- [6] R. Hirsch. *Peirce algebras and Boolean modules*, J Logic Computation **17**, 255-283 (2007).
- [7] S. M. Pinto, M.T. Oliveira-Martins, and M.C. Pinto. *Monadic Dynamic Algebras*, Math. Log. Quart. **52**, 134-150 (2006).
- [8] S. M. Pinto, M.T. Oliveira-Martins, and M.C. Pinto. *Congruences on Dynamic Algebras*, Logic J. IGPL **16**, 15-31 (2008).
- [9] A. Tarski. *On the calculus of relations*, J. Symbolic Logic **6**, 73-89 (1941).
- [10] A. Tarski. *Contributions to the theory of models, Part II*, Indagationes Mathematicae **17**, 56-64 (1955).

SANDRA MARQUES PINTO

DEPARTAMENTO DE MATEMÁTICA, UNIVERSIDADE DE COIMBRA, 3001-454 COIMBRA, PORTUGAL

E-mail address: sandra@mat.uc.pt

M. TERESA OLIVEIRA-MARTINS

DEPARTAMENTO DE MATEMÁTICA, UNIVERSIDADE DE COIMBRA, 3001-454 COIMBRA, PORTUGAL

E-mail address: meresa@mat.uc.pt